

General Reference: [Lee et al. \[2018\]](#)

## 1 Motivation

The problem we are considering is to minimise a convex function over a convex set:

$$\min_{x \in K} f(x).$$

Classical cutting-plane methods require a **SEPARATION oracle** for  $K$ : given a point outside  $K$ , find a hyperplane that separates the point from  $K$ . But what if we only have a weaker **MEMBERSHIP oracle** that can decide whether a point belongs to  $K$ ? This lecture shows how to build a separation oracle from a membership oracle, with some small residuals.

**Given:** Suppose we know a point  $x_0 \in K$  and numbers  $0 < r < R$  for the convex set  $K$  such that.

$$B(x_0, r) \subseteq K \subseteq B(x_0, R)$$

We have the access to oracle  $\text{MEM}(K)$  that determines if  $y$  is in  $K$  in the form:

**MEM**( $y \in \mathbb{R}^n, \delta > 0$ )

Output: with probability  $1 - \delta$ ,

If  $y \in B(K, \delta)$  return TRUE,

Else if  $y \notin B(K, -\delta)$  return FALSE,

where  $B(K, -\delta)$  and  $B(K, \delta)$  are defined below.

**Definition 1.** We define the “ball” around  $K$  to be the set of points at most  $\delta$  distance from  $K$ :

$$B_p(K, \delta) := \{x \in \mathbb{R}^n : \exists y \in K \text{ s.t. } \|x - y\|_p \leq \delta\}.$$

By similar ideas, we also let

$$B_p(K, -\delta) := \{x \in \mathbb{R}^n : B_p(x, \delta) \subseteq K\}.$$

**Goal:** We want to implement a  $\text{SEP}(K, x)$  to find a separating hyperplane between  $K$  and  $x$  in the form:

**SEP:** Input:  $y \in \mathbb{R}^n, \delta > 0$

Output: with prob  $1 - \delta$  :

If  $y \in B(K, \delta)$  output TRUE

Else  $c \in \mathbb{R}^n$  s.t.  $\langle c, x \rangle \leq \langle c, y \rangle + \delta \quad \forall x \in B(K, -\delta)$

The output  $c$  defines a half-space that contains most of  $K$  and puts  $y$  on the other side. For simplicity, we write  $\tilde{O}(f(n))$  for  $O(f(n))$  up to polylogarithmic factors, i.e.  $\tilde{O}(f(n)) = O(f(n) \log^k(n))$  for any  $k \in \mathbb{N}$ .

## 2 Main Results

We want to finally show the following result:

**Theorem 2** (Lemma 13 in Lee et al. [2018]). *The Algorithm 1 below gives a half space containing  $K$  with probability  $1 - \delta$ .*

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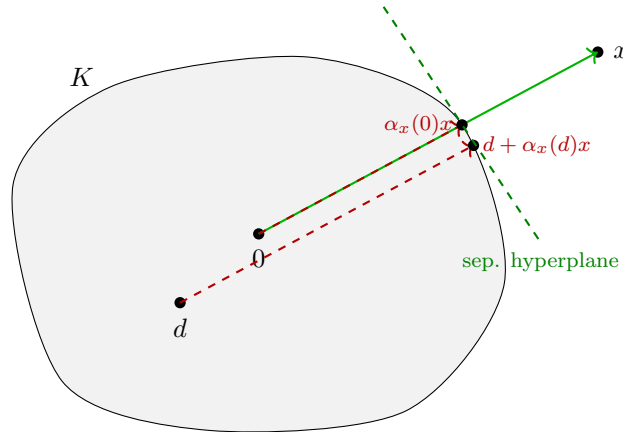
**Algorithm 1**  $\text{Separate}_{\delta, P}(K, x)$

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- 1:  $B(x_0, r) \subset K \subset B(x_0, R)$ .
- 2: **if** MEM( $K, x$ ) **then**
- 3:   return True
- 4: **else if**  $x \notin B(0, R)$
- 5:   return  $y : 0 \geq \langle y - x, x \rangle$
- 6: **end if**
- 7: Use the subgradient algorithm (Algorithm 2) to obtain an approximate subgradient  $\tilde{g}$  of  $h_x$  at 0.
- 8: **return** the half-space

$$\left\{ y : \frac{50}{\delta} n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa \geq \langle \tilde{g}, y - x \rangle \right\}$$


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**Figure 1:** Intuition for the depth function  $h_x$  and the separating half-space.

**Intuition** The core construction is a one-dimensional “depth” function in the direction of  $x$ . For any point  $d \in K$ , let

$$\alpha_x(d) := \sup\{\alpha \geq 0 \mid d + \alpha x \in K\}$$

be the largest step we can take from  $d$  in the direction  $x$  while staying inside  $K$ . Then define

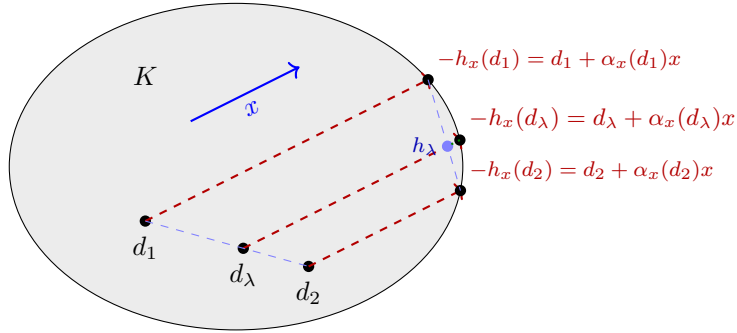
$$h_x(d) := -\alpha_x(d) \|x\|_2,$$

a non-positive function on  $K$  that measures how “deep”  $d$  is from the boundary. Since  $h_x$  is convex and  $L$ -Lipschitz as shown below, a subgradient  $g$  of  $h_x$  at 0 is bounded and satisfies  $\langle g, y \rangle \leq h_x(y) - h_x(0)$  for all  $y$ . Hence  $h_x(y) \leq 0$  on  $K$  gives the inequality that separates  $x$  from  $K$ . However, it remains as a question that how we evaluate the gradient of  $h_x$  at 0 using membership oracles, which will be discussed in Lemma 5 next section.

In the following, we will first prove the convexity and Lipschitzness of  $h_x(d)$  in Lemma 3 and 4, then give the intuition for the subgradient algorithm in Lemma 5, and finally show the main result in Section 4.

### 3 Technical Lemmas

**Lemma 3** (Lemma 11 in Lee et al. [2018]).  $h_x(d)$  is a convex function on  $K$ .



**Figure 2:** Graph for Lemma 3: The function  $-h_x(d)$  is concave.  $d_\lambda$  is a linear combination of  $d_1$  and  $d_2$ . Then  $h_\lambda$ , the linear combination of  $h_x(d_1)$  and  $h_x(d_2)$ , is on the line segment between  $d_\lambda$  and  $-h_x(d_\lambda)$ .

*Proof.* Let  $d_1, d_2 \in K$ , and  $\lambda \in [0, 1]$ . Since

$$\begin{aligned} d_1 + \alpha_x(d_1)x &\in K, \\ d_2 + \alpha_x(d_2)x &\in K, \end{aligned}$$

We have

$$[\lambda d_1 + (1 - \lambda)d_2] + [\lambda \alpha_x(d_1) + (1 - \lambda)\alpha_x(d_2)]x \in K$$

Denote  $d = (\lambda d_1 + (1 - \lambda)d_2) \in K$  and  $\beta = \lambda \alpha_x(d_1) + (1 - \lambda)\alpha_x(d_2)$ , then by the definition of  $\alpha_x$ ,  $d + \beta x \in K$  yields

$$\alpha_x(d) = \sup\{\alpha : d + \alpha x \in K\} \geq \beta.$$

Hence the definition of  $d$  and  $\beta$  gives

$$\alpha_x(\lambda d_1 + (1 - \lambda)d_2) \geq \lambda \alpha_x(d_1) + (1 - \lambda)\alpha_x(d_2)$$

Multiplying  $-\|x\|$  on both sides and using that  $h_x(d) = -\|x\|\alpha_x(d)$ , we have

$$h_x(\lambda d_1 + (1 - \lambda)d_2) \leq \lambda h_x(d_1) + (1 - \lambda)h_x(d_2),$$

which proves the convexity of  $h_x$ . ■

**Lemma 4** (Lemma 12 in Lee et al. [2018]).  $h_x$  is  $L = \frac{R+\delta}{r-\delta}$ -Lipschitz on  $B_2(0, \delta)$  for all  $\delta < r$ .

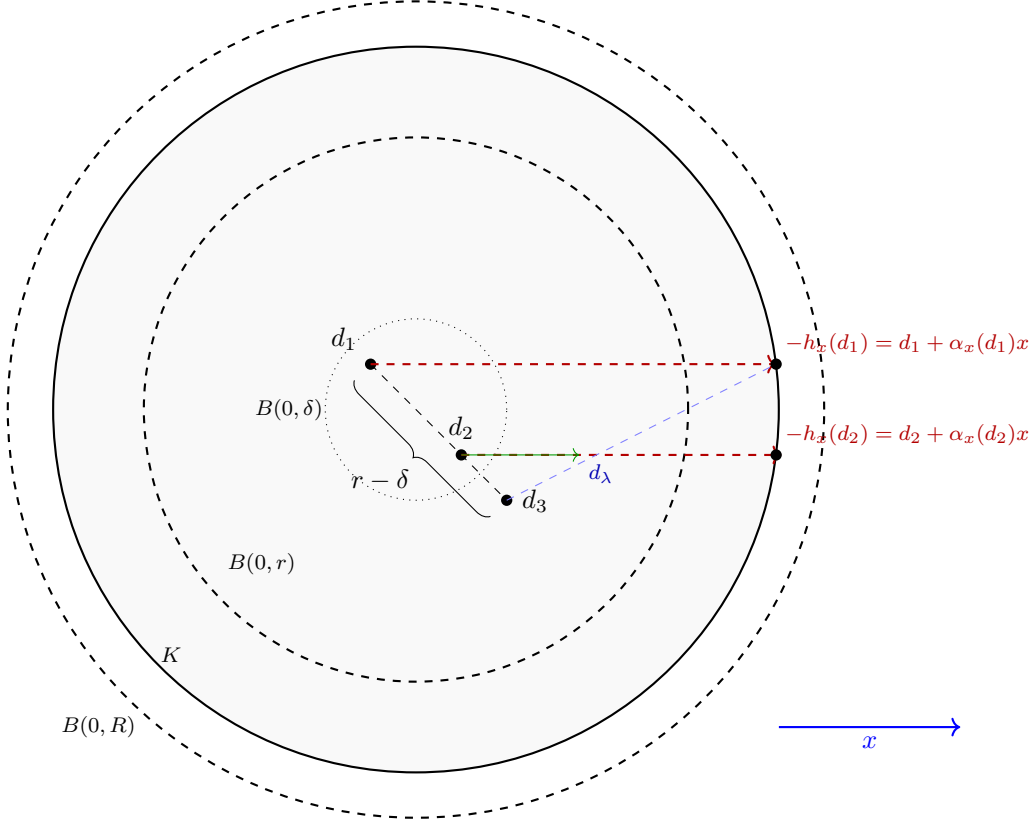
*Proof.* Let  $d_1, d_2 \in B(0, \delta)$ . Without losing of generality, we assume  $\alpha_x(d_1) \geq \alpha_x(d_2)$ , then we have

$$|h_x(d_1) - h_x(d_2)| = \|x\|_2[\alpha_x(d_1) - \alpha_x(d_2)],$$

and we therefore want to show that

$$\|x\|_2[\alpha_x(d_1) - \alpha_x(d_2)] \leq L\|d_1 - d_2\|.$$

**Case 1:**  $\|d_2 - d_1\|_2 \leq r - \delta$ .



**Figure 3:** Lemma 4, Case 1: Construct  $d_3$  by extending  $d_1d_2$  to the length of  $r - \delta$  and denote  $\lambda = \|d_1d_2\|/(r - \delta)$ . Suppose  $(-h_x(d_1))d_3$  intersects  $(-h_x(d_2))d_2$  at  $d_\lambda$ , then  $d_2d_\lambda = (1 - \lambda)\alpha_x(d_1)\|x\|$  by similar triangles. We only need to show that  $\|x\|(\alpha_x(d_1) - \alpha_x(d_2)) \leq \lambda\alpha_x(d_1)\|x\|$ , equivalently  $(1 - \lambda)\alpha_x(d_1)\|x\| \leq \alpha_x(d_2)\|x\|$ , which could be concluded by  $\|d_2d_\lambda\| \leq \|d_2(-h_x(d_2))\|$ .

Let us consider  $d_3 := d_1 + \frac{1}{\lambda}(d_2 - d_1)$  where  $\lambda = \frac{\|d_2 - d_1\|_2}{r - \delta} \in [0, 1]$ . Then we have

$$\lambda d_3 + (1 - \lambda)[d_1 + \alpha_x(d_1)x] = d_2 + (1 - \lambda)\alpha_x(d_1)x. \quad (3.1)$$

Besides, we also have

$$\|d_3\|_2 \leq \|d_1\|_2 + \frac{1}{\lambda}\|d_2 - d_1\|_2 \leq \delta + \frac{1}{\lambda}\|d_2 - d_1\|_2 = r,$$

where the first inequality comes from the Pythagorean Theorem, the second comes from  $d_1 \in B_2(0, \delta)$ , and the last comes from the definition of  $\lambda$ . Hence we have  $d_3 \in B(0, r) \subseteq K$ . Meanwhile, the definition of  $\alpha_x$  yields  $(d_1 + \alpha_x(d_1)x) \in K$ . Thus, their convex combination

$$\lambda d_3 + (1 - \lambda)[d_1 + \alpha_x(d_1)x] \in K.$$

Thus, by (3.1) and the definition of  $\alpha_x$ , we get

$$\alpha_x(d_2) = \sup\{\alpha : d_2 + \alpha x\} \geq (1 - \lambda)\alpha_x(d_1) = \left(1 - \frac{\|d_2 - d_1\|_2}{r - \delta}\right)\alpha_x(d_1).$$

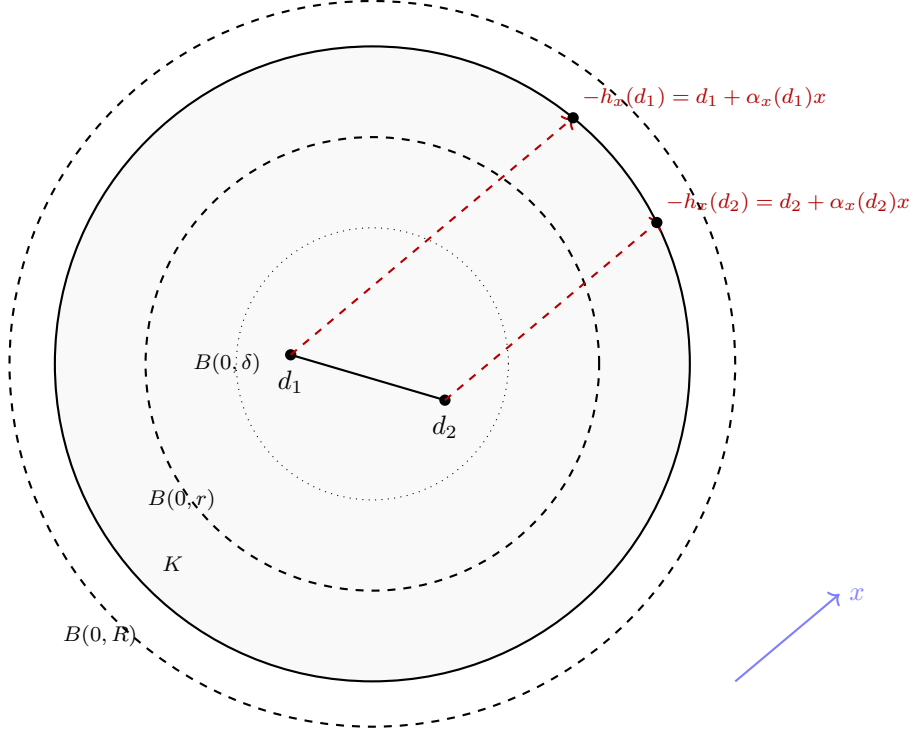
Therefore, we have

$$\|x\|_2[\alpha_x(d_1) - \alpha_x(d_2)] \leq \frac{\|d_2 - d_1\|_2}{r - \delta}\alpha_x(d_1)\|x\|_2.$$

Now we need to control  $\alpha_x(d_1)\|x\|_2$ .  $(d_1 + \alpha_x(d_1)x) \in K \subseteq B(0, R)$  and  $\|d_1\| \leq \delta$  give

$$\begin{aligned} \alpha_x(d_1)\|x\| &\leq \|d_1 + \alpha_x(d_1)x - d_1\| \\ &\leq \|d_1 + \alpha_x(d_1)x\| + \|d_1\| \\ &\leq R + \delta, \end{aligned}$$

which concludes the result.



**Figure 4:** Lemma 4, Case 2: when  $\|d_1 - d_2\| \geq r - \delta$ , the Lipschitz bound follows directly from  $\|h_x(d_1) - h_x(d_2)\| \leq R + \delta$ .

**Case 2:**  $\|d_2 - d_1\| \geq r - \delta$ . Since  $d_1, d_2 \in B(0, \delta)$ , the distance from  $d_i$  to  $B(0, R)$  is less than or equals to  $\delta + R$  for  $i = 1, 2$ . Hence we have  $h_x(d_1), h_x(d_2) \in [-R - \delta, 0]$ .

$$\begin{aligned} |h_x(d_1) - h_x(d_2)| &\leq (R + \delta) \cdot 1 \\ &\leq (R + \delta) \frac{\|d_2 - d_1\|_2}{r - \delta}, \end{aligned}$$

which gives the result. ■

With the Lipschitzness of the target function, we can now approximate the subgradient near the desire point.

**Lemma 5** (Lemma 10 in Lee et al. [2018]). *Given  $r_1, L > 0$ . Let  $f$  be a convex function on  $B_\infty(x, 2r_1)$ . Suppose that  $\|\partial f(z)\|_\infty \leq L$  for any  $z \in B_\infty(x, 2r_1)$ . Also, assume that we can compute function  $f$  with  $\varepsilon$  additive error with  $\varepsilon \leq r_1 \sqrt{n}L$ . Let  $\tilde{g}$  be an approximate subgradient function of  $f$  at  $x$  given by Algorithm 2 defined below. Then, there is a random variable  $\zeta \geq 0$  with  $\mathbb{E}\zeta \leq 3\sqrt{\frac{L\varepsilon}{r_1}}n^{5/4}$  such that*

$$f(q) \geq f(x) + \langle \tilde{g}, q - x \rangle - \zeta \|q - x\|_\infty - 4nr_1L \quad \text{for all } q \in K.$$

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**Algorithm 2** `ApproxSubGrad`( $f, x, r_1, \varepsilon$ ) (Algorithm 2: `SeparateConvexFunc` in Lee et al. [2018])

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- 1: **Require:**  $r_1 > 0$ ,  $\|\partial f(z)\|_\infty \leq L$  for any  $z \in B_\infty(x, 2r_1)$ .
- 2: Set  $r_2 = \sqrt{\frac{\varepsilon r_1}{\sqrt{n}L}}$ .
- 3: Sample  $y \in B_\infty(x, r_1)$  and  $z \in B_\infty(y, r_2)$  independently and uniformly at random.
- 4: **for**  $i = 1, 2, \dots, n$  **do**
- 5:   Let  $\alpha_i$  and  $\beta_i$  denote the end points of the interval  $B_\infty(y, r_2) \cap \{z + se_i : s \in \mathbb{R}\}$ .
- 6:   Set  $\tilde{g}_i = \frac{f(\beta_i) - f(\alpha_i)}{2r_2}$  where we compute  $f$  with  $\varepsilon$  additive error.
- 7: **end for**

**Output**  $\tilde{g}$  as the approximate subgradient of  $f$  at  $x$ .

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**Intuition** We have a convex function  $f$  (in our case  $f = h_x$ ) that is Lipschitz, and we can only obtain noisy function values by binary search on EVALUATION oracles. The algorithm first picks a random point  $y$  near the origin (inside a small ball of radius  $r_1, B(x, r_1)$ ), then approximates the partial gradients at  $y$  by taking differences on a segment in a direction of an infinite ball near  $y$ . More formally speaking:

We further randomly choose a point  $z$  in a smaller ball with radius  $r_2, B(y, r_2)$ . For each coordinate  $i$ , look at the line through  $z$  in direction  $e_i$ ; the intersection of this line with the ball  $B_\infty(y, r_2)$  is a segment  $[\alpha_i, \beta_i]$ . Estimate the  $i$ -th partial derivative as

$$\tilde{g}_i = \frac{f(\beta_i) - f(\alpha_i)}{2r_2},$$

using noisy function evaluations.

Omitting many technical details, the key outcome is:

$$f(q) \geq f(0) + \langle \tilde{g}, q \rangle - \zeta \|q\|_\infty - O(nr_1L),$$

where  $\zeta$  is a random variable with small mean. If we ignore the  $O(\cdot)$  terms and the random  $\zeta$ , we have the standard supporting hyperplane inequality for a convex function.

*Sketch of Proof.* By the convexity of  $f$ , we have for all  $q$  in  $K$ :

$$f(q) \geq f(z) + \langle \nabla f(z), q - z \rangle \tag{3.2}$$

$$= f(z) + \langle \tilde{g}, q - x \rangle + \langle \nabla f(z) - \tilde{g}, q - x \rangle + \langle \nabla f(z), x - z \rangle. \tag{3.3}$$

Now we have  $f(z) \geq f(x) - L\|z - x\|$  by L-Lipschitz,  $\langle \nabla f(z), x - z \rangle$  lower bounded by  $-\|\nabla f(z)\|_\infty \|x - z\|_1$  thus also  $O(L\|z - x\|)$  by Lipschitz. They are both  $O(nr_1L)$  because  $z$  is selected from the small ball near  $x$ . Since  $\langle \nabla f(z) - \tilde{g}, q - x \rangle \geq -\|\nabla f(z) - \tilde{g}\|_1 \|q - x\|_\infty$ , it remains to show that the random variable  $\|\nabla f(z) - \tilde{g}\|_1$  has a small mean, which is proven by Lemma 6 below. ■

**Lemma 6.** For any  $0 < r_2 \leq r_1$  and twice differentiable convex function  $f$  defined on  $B_\infty(x, r_1 + r_2)$  with  $\|\nabla f(z)\|_\infty \leq L$  for any  $z \in B_\infty(x, r_1 + r_2)$  we have

$$\mathbb{E}_{y \in B_\infty(x, r_1)} \mathbb{E}_{z \in B_\infty(y, r_2)} \|\nabla f(z) - g(y)\|_1 \leq n^{3/2} \frac{r_2}{r_1} L$$

where  $g(y)$  is the average of  $\nabla f$  over  $B_\infty(y, r_2)$ .

*Sketch of Proof.* The idea is that  $\|\nabla f(z) - \tilde{g}\|_1$  is bounded by the second derivative of  $f$  at  $z$ . By Stoke's Theorem, given the expectation integrates the second derivative, the term is controlled by integration of  $\langle \nabla f, n \rangle$  over the boundary, where  $n$  is the normal vector. Therefore,  $\mathbb{E}[\|\nabla f(z) - \tilde{g}\|_1]$  is bounded. ■

## 4 Proof of Main Results

Now we begin the proof of the main theorem.

*Proof of Theorem 2.* The proof of some technical lemmas are in the next section. When  $x \notin B(0, R)$ , Algorithm 1 will output SEPARATED for  $B(0, R)$ . Thus in the rest of our proof, we may assume  $x \in B(0, R)$  but  $x \notin B(K, -\varepsilon)$ .

For simplicity, we denote  $\kappa = \frac{R}{r}$ . By Lemma 3,  $h_x$  is a convex function. Moreover, given  $S = r/2$ , Lemma 4 yields  $h_x$  is a  $L$ -Lipschitz function with

$$L = \frac{R + r/2}{r/2} < \frac{3R/2}{r/2} = 3\kappa.$$

**Step 1: Approximate subgradient inequality.** We want to apply Lemma 5 to the function  $f = h_x$  over the domain  $K$ . Given the Lipschitz continuity, the lemma requires a small radius  $r_1$  such that  $B_\infty(0, 2r_1) \subset B_2(0, S)$ . We will choose  $r_1$  later (in terms of  $\varepsilon$ ). The lemma then supplies a random vector  $\tilde{g}$  and a random variable  $\xi \geq 0$  with

$$\mathbb{E}[\xi] \leq 3\sqrt{\frac{L\varepsilon}{r_1}} n^{5/4},$$

such that for every  $y \in K$ ,

$$h_x(y) \geq h_x(0) + \langle \tilde{g}, y \rangle - \xi \|y\|_\infty - Cnr_1L, \quad (4.1)$$

where  $C$  is a constant (the proof gives  $C = 24$ , but we keep the constant for clarity).

To approximate a separating hyperplane, we need to show  $\langle \tilde{g}, x \rangle \geq \langle \tilde{g}, y \rangle + C'$  for some constant  $C'$  independent of  $y$ . Given  $h_x(y) \leq 0$ , it remains to control  $-h_x(0) - \langle \tilde{g}, x \rangle$ .

**Step 2: Lower bound  $\langle \tilde{g}, x \rangle$  and  $h_x(0)$ .** Since  $\|\frac{x}{\kappa}\| = r \cdot \frac{\|x\|}{R} \leq r$ , we have  $\frac{-x}{\kappa} \in B(0, r) \subset K$  in the domain of  $h_x$ , therefore we can derive  $h_x\left(\frac{-x}{\kappa}\right)$  from definition:

$$\begin{aligned} h_x\left(\frac{-x}{\kappa}\right) &= -\|x\|_2 \sup \left\{ \alpha : -\frac{x}{\kappa} + \alpha x \in K \right\} \\ &= -\|x\|_2 \sup \left\{ \beta + \frac{1}{\kappa} : 0 + \beta x \in K \right\} \quad \text{by letting } \alpha - \frac{1}{\kappa} = \beta \\ &= -\|x\|_2 \left[ \alpha_x(0) + \frac{1}{\kappa} \right] \\ &= h_x(0) - \frac{\|x\|_2}{\kappa}. \end{aligned}$$

On the other hand, letting  $y = -x/\kappa$  in (4.1) yields

$$h_x\left(\frac{-x}{\kappa}\right) \geq h_x(0) + \left\langle \tilde{g}, \frac{-x}{\kappa} \right\rangle - \frac{1}{\kappa} \xi \|x\|_\infty - 12nr_1\kappa$$

Thus we can get

$$\langle \tilde{g}, x \rangle \geq \|x\|_2 - \xi \|x\|_\infty - 12nr_1\kappa^2 \quad (4.2)$$

Moreover, by Lemma 7, we obtain

$$h_x(0) \geq -\|x\|_2 - \varepsilon\kappa. \quad (4.3)$$

**Final Step: Derive separation inequality.** Combining (4.1), (4.2), and (4.3) gives

$$\begin{aligned} h_x(y) &\geq \langle \tilde{g}, y - x \rangle - \xi \|y\|_\infty - \xi \|x\|_\infty - 12nr_1\kappa - 12nr_1\kappa^2 - \varepsilon\kappa \\ &\geq \langle \tilde{g}, y - x \rangle - 2\xi R - 24nr_1\kappa^2 - \varepsilon\kappa. \end{aligned}$$

for any  $y \in K$ . Then  $h_x(y) \leq 0$  yields

$$\langle \tilde{g}, y - x \rangle \leq 2\xi R + 24nr_1\kappa^2 + \varepsilon\kappa,$$

where we define  $\tilde{\xi}$  to be the right hand side. Thus, we have  $\tilde{\xi} \geq \langle \tilde{g}, y - x \rangle$  where  $\tilde{\xi}$  is a random variable independent of  $y$ . Now we can choose  $r_1 = n^{1/6} R^{2/3} \varepsilon^{1/3} \kappa^{-1}$  and use  $\mathbb{E}\tilde{\xi} \leq 3\sqrt{\frac{12\kappa\varepsilon}{r_1}} n^{5/4}$  in Lemma 5 to have

$$\begin{aligned} \mathbb{E}\tilde{\xi} &\leq 6\sqrt{\frac{12\kappa\varepsilon}{r_1}} n^{5/4} R + 24nr_1\kappa^2 + \varepsilon\kappa \\ &\leq 45n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa + \varepsilon\kappa \\ &\leq 50n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa. \end{aligned}$$

Then the Algorithm 1 would output the half-space

$$\left\{ y : \frac{50}{\delta} n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa \geq \langle \tilde{g}, y - x \rangle \right\}$$

with probability, by Markov's inequality,

$$1 - \mathbb{P}\left(\langle \tilde{g}, y - x \rangle \geq \frac{50}{\delta} n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa\right) \geq 1 - \mathbb{P}\left(\tilde{\xi} \geq \frac{50}{\delta} n^{7/6} R^{2/3} \varepsilon^{1/3} \kappa\right) \geq 1 - \delta.$$

■

**Lemma 7.**  $h_x(0) \geq -\|x\|_2 - \varepsilon\kappa$ .<sup>1</sup>

*Proof.* For  $h_x(0)$ , we want to control  $\alpha_x(0)$ , i.e. we want to find some point out of  $K$  to cover  $\alpha_x(0)$ . In the following, we show that for all  $x \notin B(K, -\varepsilon)$ ,  $x/(1 - \frac{\varepsilon}{r}) \notin K$ .

Suppose for contradiction, there exists  $u \in K$  such that  $(1 - \frac{\varepsilon}{r})u = x$ . Then for any  $\|z - (1 - \frac{\varepsilon}{r})u\| \leq \varepsilon$ , we denote  $z := (1 - \frac{\varepsilon}{r})u + \frac{\varepsilon}{r}v$ , where  $\|v\| \leq r$ .  $B(0, r) \subset K$ , we have

$$v \in B(0, r) \subset K,$$

that is to say,  $z$  is a convex combination of  $u$  and  $v$ , both in  $K$ . Hence  $z$  is in  $K$  too. As a result, we get

$$B(x, \varepsilon) = B\left(\left(1 - \frac{\varepsilon}{r}\right)y, \varepsilon\right) \subset K.$$

which means  $x \in B(K, -\varepsilon)$ , contradiction. Given  $x/(1 - \frac{\varepsilon}{r}) \notin K$ , we have

$$0 < \alpha_x(0) \leq \left(1 - \frac{\varepsilon}{r}\right).$$

Rearranging yields

$$1 - \frac{1}{\alpha_x(0)} \leq \frac{\varepsilon}{r}$$

Multiplying  $R$  on both sides gives

$$R \left( \frac{\alpha_x(0) - 1}{\alpha_x(0)} \right) \leq \varepsilon\kappa.$$

<sup>1</sup>The authors made a typo when proving the bound but it doesn't affect the main result.

For all  $\alpha$  such that  $\alpha x \in K \subseteq B(0, R)$ , given  $\|\alpha\|\|x\| \leq R$ , we have  $\alpha_x(0)\|x\| = \|\alpha_x(0)\|\|x\| \leq R$  by the definition of  $\alpha_x(0)$ . Hence we can further extend the inequality as

$$\|x\| (\alpha_x(0) - 1) \leq \varepsilon\kappa.$$

Rearranging gives

$$h_x(0) = -\alpha_x(0)x \geq -\|x\|_2 - \varepsilon\kappa.$$

■

## References

Yin Tat Lee, Aaron Sidford, and Santosh S Vempala. Efficient convex optimization with membership oracles. In *Conference On Learning Theory*, pages 1292–1294. PMLR, 2018. [1](#), [2](#), [3](#), [5](#), [6](#)