

Convex Optimization: Problem Set 1

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Some problems borrowed from Optimization course of Daniel Dadush.

1. Recall the definition of convex hull:

$$\text{conv}(S) := \left\{ \sum_{i=1}^N \lambda_i \mathbf{x}_i \mid \mathbf{x}_i \in S, \lambda_i \geq 0, \sum_{i=1}^N \lambda_i = 1 \right\}.$$

1. Prove that for finite S , $\text{conv}(S) = \bigcap_{K \supseteq S} K$ where the intersection is over all closed convex K containing S . Therefore $\text{conv}(S)$ is the smallest convex set containing S .
2. Prove Jensen's inequality: for convex $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and input $\mathbf{x} = \sum_{i=1}^N \lambda_i \mathbf{x}_i$ where $\lambda_i \geq 0, \sum_{i=1}^N \lambda_i = 1$

$$f(\mathbf{x}) \leq \sum_{i=1}^N \lambda_i f(\mathbf{x}_i).$$

3. Show that if $\mathcal{C} \subseteq \mathbb{R}^n$ is a compact convex set and $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is a convex function then the **supremum** of f over \mathcal{C} is attained at an extreme point of \mathcal{C} .

Solution:

1. Clearly any convex $K \supseteq S$ contains any convex combination $\sum_{i=1}^N \lambda_i \mathbf{x}_i$ where $\mathbf{x}_i \in S$, so $\text{conv}(S) \subseteq \bigcap_{K \supseteq S} K$. Conversely, if $x \notin \text{conv}(S)$ then there exists a separating hyperplane $\langle w, x \rangle > \nu := \sup_{y \in \text{conv}(S)} \langle w, y \rangle$ and so the halfspace

$$H := \{y \mid \langle y, w \rangle \leq \nu\}$$

is a convex set satisfying $H \supseteq S, x \notin H$.

2. Recall that Jensen's inequality states that for convex function $f : \mathcal{C} \rightarrow \mathbb{R}$ on convex set \mathcal{C} , set of points $\mathbf{x}_1, \dots, \mathbf{x}_k \in \mathcal{C}$, and coefficients $\lambda \in \mathbb{R}_+^k$ with $\sum_{i=1}^k \lambda_i = 1$, we have

$$f\left(\sum_{i=1}^k \lambda_i \mathbf{x}_i\right) \leq \sum_{i=1}^k \lambda_i f(\mathbf{x}_i).$$

We will prove the statement by induction on the number of points k . Notice that the case $k = 2$ is exactly the definition of convexity for f . Now for $k \geq 2$, given $k + 1$ points $\mathbf{x}_1, \dots, \mathbf{x}_{k+1}$ and non-zero coefficients $\lambda \in \mathbb{R}_{++}^{k+1}$ with $\sum_{i=1}^{k+1} \lambda_i = 1$, let $\tilde{\lambda} := \sum_{j=1}^k \lambda_j = 1 - \lambda_{k+1}$ and

$$\mathbf{x} := \sum_{i=1}^k \frac{\lambda_i}{\tilde{\lambda}} \mathbf{x}_i.$$

Note this is a convex combination of $\mathbf{x}_1, \dots, \mathbf{x}_k$ as $\gamma_i := \frac{\lambda_i}{\tilde{\lambda}}$ satisfies $\gamma \in \mathbb{R}_+^k$ and $\sum_{i=1}^k \gamma_i = 1$. Therefore we can bound the function as

$$\begin{aligned} f\left(\sum_{i=1}^{k+1} \lambda_i \mathbf{x}_i\right) &= f(\tilde{\lambda} \mathbf{x} + (1 - \tilde{\lambda}) \mathbf{x}_{k+1}) \leq \tilde{\lambda} f(\mathbf{x}) + (1 - \tilde{\lambda}) f(\mathbf{x}_{k+1}) \\ &= \tilde{\lambda} f\left(\sum_{i=1}^k \frac{\lambda_i}{\tilde{\lambda}} \mathbf{x}_i\right) + \lambda_{k+1} f(\mathbf{x}_{k+1}) \leq \sum_{i=1}^k \lambda_i f(\mathbf{x}_i) + \lambda_{k+1} f(\mathbf{x}_{k+1}), \end{aligned}$$

where in the first step we used the definitions of $\mathbf{x}, \tilde{\lambda}$, in the second step we used the base case, and in the final step we used the induction hypothesis.

3. We focus on the case $\mathcal{C} = \text{conv}\{x_1, \dots, x_N\}$ and show the maximum is achieved at some x_i . By continuity of f and compactness of \mathcal{C} , $x^* := \arg \max_{x \in \mathcal{C}} f(x)$ is attained at some point. Now let $x^* = \sum_{i=1}^N \lambda_i x_i$ where $\lambda \geq 0, \sum_{i=1}^N \lambda_i = 1$. Then by Jensen's inequality we have

$$f(x^*) \leq \sum_{i=1}^N \lambda_i f(x_i) \leq \max_{i \in [N]} f(x_i),$$

where the first step was by Jensen's inequality, in the second we used that λ is a convex combination. Since $x^* := \arg \max_{x \in \mathcal{C}} f(x)$ is a maximizer, the statement is shown.

We sketch some remarks for removing technical assumptions: it can be shown that convex f is always continuous at any $x \in \text{int}(\text{dom}(f))$ (see Theorem 10.1 and Cor 10.1.1 in Rockafellar). Since our function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ never takes unbounded values, the domain is \mathbb{R}^n , so f is automatically continuous on \mathcal{C} .

For general compact \mathcal{C} , the Krein-Milman theorem tells us that \mathcal{C} is the convex hull of its extreme points, i.e. for the optimizer x^* , there exists extreme points $\{x_1, \dots, x_N\}$ and $\lambda \geq 0, \sum_{i=1}^N \lambda_i = 1$ such that $x^* = \sum_{i=1}^N \lambda_i x_i$.

2. Describe the set of boundary points for the following norm balls, i.e. points not in the interior. For each boundary point \mathbf{x} give the set of supporting hyperplanes at \mathbf{x} .

1. $B_1^n := \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{x}\|_1 := \sum_{i=1}^n |x_i| \leq 1\}$
2. $B_2^n := \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{x}\|_2 := \sqrt{\sum_{i=1}^n x_i^2} \leq 1\}$
3. $B_\infty^n := \{\mathbf{x} \in \mathbb{R}^n : \|\mathbf{x}\|_\infty := \max_{i \in [n]} |x_i| \leq 1\}$.

Solution:

1. x is on the boundary iff $\|x\|_1 = 1$, and in this case the set of supporting hyperplanes have normal vectors

$$\{w \in [-1, 1] \mid \forall i \in \text{supp}(x) : w_i := \text{sign}(x_i)\}.$$

2. x is on the boundary iff $\|x\|_2 = 1$, and in this case the supporting hyperplane has normal vector x .

3. x is on the boundary iff $\|x\|_\infty = 1$, and in this case the set of supporting hyperplanes have normal vectors

$$\text{conv}\{e_i \mid |x_i| = 1\}.$$

3. For convex closed $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and $x \in \text{dom}(f)$:

- Show for any subgradient $g \in \partial f(x)$, $(g, -1)$ gives a *supporting hyperplane* at $(x, f(x))$ for the epigraph

$$\text{epi}(f) := \{(x, t) \mid f(x) \leq t\}.$$

- Show that $g \in \partial f(x)$ gives a *supporting hyperplane* at x for the sub-level set

$$L := \{y \in \mathbb{R}^n \mid f(y) \leq f(x)\}.$$

Solution:

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$$\begin{aligned} g \in \partial f(x) &\iff \forall y : f(y) \geq f(x) + \langle g, y - x \rangle \\ &\iff \langle (g, -1), (x, f(x)) \rangle \geq \sup_y \langle (g, -1), (y, f(y)) \rangle = \sup_{(y,t) \in \text{epi}(f)} \langle (g, -1), (y, t) \rangle \end{aligned}$$

where the first step was by definition, and in the last step we used $(y, t) \in \text{epi}(f) \iff f(y) \leq t$.

- We claim $L \subseteq \{y \mid \langle g, y \rangle \leq \langle g, x \rangle\} =: H$, which is clearly a supporting hyperplane at x . To show this, assume $y \notin H$ so $\langle g, y \rangle > \langle g, x \rangle$. Then by definition of subgradient we have

$$f(y) - f(x) \geq \langle g, y - x \rangle > 0,$$

i.e. $y \notin L$ the sub-level set.

4. Prove Exercise 2.27 in [BV]: Let $K \subseteq \mathbb{R}^n$ be closed, bounded, with non-empty interior, such that there exists a supporting hyperplane of K at every point of the boundary $x \in \partial K$. Show this implies K is convex.

Solution:

[Better proof from Student Solutions]:

We show the assumption implies $K = \bigcap_{H \supseteq K} H =: C$, where the intersection is over all supporting halfspaces at boundary points. The containment $K \subseteq C$ is clear since each halfspace contains K . For the reverse, assume for contradiction there is some $x \in C, x \notin K$. Consider the line segment $[x, x_0]$ to $x_0 \in \text{int}(K)$. This segment intersects the boundary

$$\exists \lambda \in (0, 1) : x_\lambda := (1 - \lambda)x_0 + \lambda x \in \partial K.$$

Let $v := x - x_0$ so $x_\lambda = x_0 + \lambda v$. By the assumption, there is a supporting hyperplane

$$\langle w, x_\lambda \rangle \geq \sup_{z \in K} \langle w, z \rangle.$$

Since x_0 is in the interior, we must have $\langle w, x \rangle > \langle w, x_0 \rangle$, which then implies $\langle w, x \rangle > \langle w, x_\lambda \rangle$. But then this implies $x \notin C$.

5. Affine and quadratic functions are the most basic convex functions. We will prove some properties about them:

- Show $h : \mathbb{R}^n \rightarrow \mathbb{R}$ is convex *and* concave (i.e. h is convex and $-h$ is convex) iff h is an affine function, i.e. $h(x) = \langle a, x \rangle + b$ for $a \in \mathbb{R}^n, b \in \mathbb{R}$.
- Let $\tilde{q}(x) := \langle x, Qx \rangle + \langle a, x \rangle + b$ for symmetric matrix $Q \in \mathbb{R}^{n \times n}, a \in \mathbb{R}^n, b \in \mathbb{R}$. Show \tilde{q} is convex iff $q(x) := \langle x, Qx \rangle$ is convex. (Hint: use first part)
- Show $q(x) = \langle x, Qx \rangle$ is convex iff

$$\forall v \in \mathbb{R}^n : \langle v, Qv \rangle \geq 0,$$

and similarly show it is strictly convex iff the inequality is strict for any $v \neq 0$. (These conditions are known as positive-semi-definiteness and positive-definiteness, and are denoted $Q \succeq 0, Q \succ 0$.)

- Find the optimizer and optimum value of strictly convex quadratic

$$\tilde{q}(x) := \langle x, Qx \rangle + \langle a, x \rangle + b.$$

Solution:

- Clearly affine h is convex and concave, so we focus on the converse. By convexity and concavity, respectively, we have

$$\lambda h(x) + (1 - \lambda)h(y) \leq h(\lambda x + (1 - \lambda)y) \leq \lambda h(x) + (1 - \lambda)h(y),$$

so in fact we must have equality for all $x, y \in \mathbb{R}^n, \lambda \in [0, 1]$. We can use this to explicitly compute the affine function:

$$h(x) = h(0) + \sum_{i=1}^n (h(e_i) - h(0))x_i.$$

- We use that sums of convex functions are convex, and

$$\tilde{q}(x) = q(x) + h(x), \quad q(x) = \tilde{q}(x) - h(x)$$

where $h(x) := \langle a, x \rangle + b$ is an affine function, so convex and concave. So \tilde{q} is convex iff q is.

- We first note that

$$\begin{aligned} \lambda x x^T + (1 - \lambda) y y^T - (\lambda x + (1 - \lambda) y)(\lambda x + (1 - \lambda) y)^T \\ = (\lambda - \lambda^2) x x^T + ((1 - \lambda) - (1 - \lambda)^2) y y^T - \lambda(1 - \lambda)(x y^T + y x^T) \\ = \lambda(1 - \lambda)[x x^T + y y^T - x y^T - y x^T] = \lambda(1 - \lambda)(x - y)(x - y)^T. \end{aligned}$$

Therefore we can rewrite

$$\begin{aligned} \lambda q(x) + (1 - \lambda)q(y) - q(\lambda x + (1 - \lambda)y) \\ = \langle Q, \lambda x x^T + (1 - \lambda) y y^T - (\lambda x + (1 - \lambda) y)(\lambda x + (1 - \lambda) y)^T \rangle \\ = \lambda(1 - \lambda) \langle Q, (x - y)(x - y)^T \rangle. \end{aligned}$$

Now q is convex iff the above is non-negative for all $x, y \in \mathbb{R}^n, \lambda \in [0, 1]$, which is equivalent to the condition $\langle v, Qv \rangle \geq 0$ for $v = x - y$.

- We compute the gradient and check where it vanishes:

$$0 = \nabla \tilde{q}(x^*) = 2Qx + a \quad \iff \quad x^* = -\frac{Q^{-1}a}{2}$$

where we used that \tilde{q} is strictly convex so Q is invertible. Substituting gives

$$\tilde{q}(x^*) = \frac{\langle Q^{-1}a, Q(Q^{-1}a) \rangle}{4} - \frac{\langle a, Q^{-1}a \rangle}{2} + b = -\frac{\langle a, Q^{-1}a \rangle}{4} + b.$$

6. An Ellipsoid is an affine image of the Euclidean ball

$$\mathcal{E} = c + AB_2^n \quad \text{where} \quad B_2^n := \{x \in \mathbb{R}^n \mid \|x\|_2 \leq 1\},$$

for some $c \in \mathbb{R}^n$ and $A \in \mathbb{R}^{n \times n}$ invertible.

- Let $q(x) := \langle x, Qx \rangle + \langle d, x \rangle + e$ be *strictly convex*. Show any sub-level set

$$L_t := \{x \in \mathbb{R}^n \mid q(x) \leq t\}$$

is either empty or an Ellipsoid (i.e. find c, A such that $L_t = c + AB_2$).

- Conversely, given Ellipsoid $\mathcal{E} = c + AB_2$ as above, find convex quadratic $q(x) := \langle x, Qx \rangle + \langle d, x \rangle + e$ such that

$$\mathcal{E} = \{x \in \mathbb{R}^n \mid q(x) \leq 1\}.$$

Solution:

- By the last part of the above question, we can compute the optimizer

$$x^* := \arg \min_x q(x) = -\frac{Q^{-1}d}{2} \quad \text{with} \quad q(x^*) = e - \frac{\langle d, Q^{-1}d \rangle}{4}.$$

so clearly L_t is empty for $t < q(x^*)$. We can rewrite q in terms of the optimizer

$$q(x) = \langle (x - x^*), Q(x - x^*) \rangle + q(x^*),$$

as can be verified directly. Therefore, for any $t \geq q(x^*)$ we have

$$\begin{aligned} x \in L_t &\iff \langle (x - x^*), Q(x - x^*) \rangle \leq t - q(x^*) \\ &\iff \|Q^{1/2}(x - x^*)\|_2 \leq \sqrt{t - q(x^*)} \\ &\iff x - x^* \in \sqrt{t - q(x^*)} Q^{-1/2} B_2^n, \end{aligned}$$

where in the last step we used that Q is invertible. This gives the required form for the Ellipsoid.

- We reverse the sequence of equivalences

$$\begin{aligned} x \in \mathcal{E} = c + AB_2^n &\iff x - c \in AB_2^n \iff A^{-1}(x - c) \in B_2^n \\ &\iff 1 \geq \|A^{-1}(x - c)\|_2^2 = \langle (x - c), (AA^T)^{-1}(x - c) \rangle, \end{aligned}$$

so $q(x) := \langle (x - c), (AA^T)^{-1}(x - c) \rangle$ is the required quadratic.

7. Recall the GLS oracle model for convex sets. Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be a convex closed function with $\text{dom}(f) = \mathbb{R}^n$. We investigate natural function oracles in terms of the epigraph:

- Show $\text{MEM}(\text{epi}(f))$ can be implemented using an EVALuation oracle for f ; show $\text{EVAL}(f)$ can be approximately implemented using $\text{MEM}(\text{epi}(f))$, i.e. for input $x \in \mathbb{R}^n$ compute $t \in \mathbb{R}$ such that $|t - f(x)| \leq \varepsilon$.
- Show $\text{SEP}(\text{epi}(f))$ can be implemented using an GRADient (or subgradient) and EVALuation oracles for f ; show $\text{GRAD}(f)$ can be approximately implemented using $\text{SEP}(\text{epi}(f))$, i.e. for input $x \in \mathbb{R}^n$ compute $g \in \mathbb{R}^n$ such that

$$\forall y \in \mathbb{R}^n : f(y) \geq f(x) - \varepsilon + \langle g, y - x \rangle.$$

- Relate $\text{OPT}(\text{epi}(f))$ and EVAL and GRAD for the Fenchel dual

$$f^*(w) := \sup_{x \in \mathbb{R}^n} \langle w, x \rangle - f(x).$$

- Relate $\text{MEM}(\text{epi}(f^*))$ and VAL($\text{epi}(f)$).

Solution:

- $(x, t) \in \text{epi}(f) \iff f(x) \leq t$, so for query $\text{MEM}(\text{epi}(f), (x, t))$ we can compute $\text{EVAL}(f, x)$ and output YES iff $f(x) \leq t$. For the opposite direction, we can apply binary search on t to find an ε -approximation for $f(x)$ in $\log \frac{f_{\max} - f_{\min}}{\varepsilon}$ time, where in order for this to be finite, we need the assumption that we have some bounds $f_{\min} \leq f(x) \leq f_{\max}$.
- The membership part of $\text{SEP}(\text{epi})$ is implemented using the EVALuation oracle for f as above. So assume query $(x, t) \notin \text{epi}(f)$. Then by part (1) of question 3 above, we have that for any $g \in \partial f(x)$ $(g, -1)$ is a supporting hyperplane for the epigraph at $(x, f(x))$, so we can use the GRADient oracle to output this separating hyperplane.

For the reverse direction, we can first use our SEPARation oracle to compute t such that $f(x) > t \geq f(x) - \varepsilon$ using binary search. Further, $(x, t) \notin \text{epi}(f)$ so the oracle outputs separating hyperplane $(\tilde{g}, -\nu)$ such that

$$\langle (\tilde{g}, -\nu), (x, t) \rangle = \langle \tilde{g}, x \rangle - \nu t > \sup_{(y, s) \in \text{epi}(f)} \langle (g, -\nu), (y, s) \rangle.$$

Since $\text{epi}(f)$ is upwards closed, we must have $\nu \geq 0$. Further, since $\text{dom}(f) = \mathbb{R}^n$, we must in fact have $\nu > 0$. Therefore, by normalizing $g := \tilde{g}/\nu$ we have

$$\langle g, x \rangle - t > \sup_y \langle g, y \rangle - f(y),$$

so by rearranging and using that $f(x) > t \geq f(x) - \varepsilon$ we have that g is an approximate subgradient.

- $f^*(w) = h_{\text{epi}(f)}(w, -1)$ so the value output by $\text{OPT}(\text{epi}(f), (w, -1))$ is equivalent to $\text{EVAL}(f^*, w)$. Further it can be directly verified that

$$x \in \partial f^*(w) \iff f^*(w) = \langle w, x \rangle - f(x),$$

so the optimizer output by $\text{OPT}(\text{epi}(f), (w, -1))$ is equivalent to $\text{GRAD}(f^*, w)$.

- $(w, \nu) \in \text{epi}(f^*) \iff f^*(w) \leq \nu \iff h_{\text{epi}(f)}(w, -1) \leq \nu \iff \langle (w, -1), \cdot \rangle \leq \nu$ if VALid for $\text{epi}(f)$.